#### Which Ball is Heavier?

# **COMP108 Algorithmic Foundations**

**Mathematical Induction** 

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# balance scale 9 balls look identically the same but 1 is heavier than the rest How to find the heavier one

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#### Learning outcomes

- > Understand the concept of Induction
- > Able to prove by Induction

#### **Analysis of Algorithms**

After designing an algorithm, we analyze it.

> Proof of correctness: show that the algorithm gives the desired result

by weighing 2 times only?

- > Time complexity analysis: find out how fast the algorithm runs
- > Space complexity analysis: find out how much memory space the algorithm requires
- > Look for improvement: can we improve the algorithm to run faster or use less memory? is it best possible?

(Induction)

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(Induction)

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#### A typical analysis technique

#### Induction

> technique to prove that a property holds for all natural numbers (or for all members of an infinite sequence)

E.g., To prove  $1+2+...+n = n(n+1)/2 \forall +ve integers n$ 

` '			
n	LHS	RHS	LHS = RHS?
1	1	1*2/2 = 1	
2	1+2 = 3	2*3/2 = 3	
3	1+2+3 = 6	3*4/2 = 6	

However, this isn't a proof and we cannot enumerate over all possible numbers.

 $\Rightarrow$  Induction

#### Intuition - Long Row of Dominoes

- > How can we be sure each domino will fall?
- > Enough to ensure the 1st domino will fall?
  - > No. Two dominoes somewhere may not be spaced properly







- > Enough to ensure all are spaced properly?
  - > No. We need the 1st to fall
- > Both conditions required:
  - > 1st will fall; & after the kth fall, k+1st will also fall
  - > then even infinitely long, all will fall



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#### Induction

To prove that a property holds for every positive integer n

Two steps

- > Base case: Prove that the property holds for n = 1
- > Induction step: Prove that if the property holds for n = k (for some positive integer k), then the property holds for n = k + 1
- > Conclusion: The property holds for every positive integer n

#### Example

To prove:  $1 + 2 + 3 + ... + n = \frac{n(n+1)}{2}$  \(\forall + ve \) integers n

- > Base case: When n=1 LHS=1, RHS=  $\frac{1\times 2}{2}$  = 1. So, the property holds for n=1.
- > Induction hypothesis:

Assume that the property holds when n=k for some

- i.e., assume that  $1 + 2 + 3 + ... + k = \frac{k(k+1)}{2}$
- > Induction step: When n=k+1, LHS becomes 1 + 2 + 3 + ... + k + (k+1)RHS becomes  $\frac{(k+1)((k+1)+1)}{(k+1)(k+1)}$
- .. we want to prove

 $1 + 2 + 3 + \dots + k + (k+1) = \frac{(k+1)(k+2)}{2}$ 

(Induction)

Target: to prove
$$1 + 2 + ... + k + (k + 1) = \frac{(k + 1)(k + 2)}{2}$$

#### Induction Step: When n=k+1

- > So, property also holds for n=k+1
- > Conclusion: property holds for all +ve integers n

(Induction)

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#### Example 2

To prove  $\underline{n^3+2n}$  is divisible by 3  $\forall$  integers  $n\geq 1$ 

n	n³+2n	divisible by 3?
1	1+2 = 3	
2	8+4 = 12	
3	27+6 = 33	
4	64+8 = 72	©

Prove it by induction...

(Induction)

Induction hypothesis k<sup>3</sup>+2k is divisible by 3

Target: to prove  $(k+1)^3+2(k+1)$  is divisible by 3

>Induction step: When n=k+1,

$$(k+1)^{3}+2(k+1) = (k^{2}+2k+1)(k+1) + (2k+2)$$

$$= (k^{3}+3k^{2}+3k+1) + (2k+2)$$

$$= (k^{3}+2k) + 3(k^{2}+k+1)$$
sum is divisible by 3

by hypothesis, divisible by 3 divisible by 3

- >Property holds for n=k+1
- >By principle of induction: property holds ∀ integers n≥1

(Induction

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### Example 3

To prove  $2^n < n! \forall +ve \text{ integers } n \ge 4$ .

- Base case: When n=4,
   LHS = 2<sup>4</sup> = 16, RHS = 4! = 4\*3\*2\*1 = 24,
   LHS < RHS</li>
   So, property holds for n=4
- ➤ Induction hypothesis: Assume property holds for n=k for some integer k ≥ 4, i.e., assume 2<sup>k</sup> < k!</p>
- ➤ Induction step: When n=k+1, LHS becomes 2<sup>k+1</sup>, RHS becomes (k+1)!

Target: to prove  $2^{k+1} < (k+1)!$ 

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#### Conclusion

#### We have proved

- 1. property holds for n=1
- 2. if property holds for n=k, then also holds for n=k+1

#### In other words,

- > holds for n=1 implies holds for n=2 (induction step)
- > holds for n=2 implies holds for n=3 (induction step)
- > holds for n=3 implies holds for n=4 (induction step)
- > and so on .....

By principle of induction: holds for all +ve integers n

(Induction)

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#### Example 2

To prove  $n^3+2n$  is divisible by 3  $\forall$  integers  $n\geq 1$ 

- Base case: When n=1, n³+2n=1+2=3, divisible by 3. So property holds for n=1.
- Induction hypothesis: Assume property holds for n=k, for some +ve int k, i.e., assume k³+2k is divisible by 3
- Induction step: When n=k+1, LHS becomes (k+1)<sup>3</sup> + 2(k+1)

Target: to prove (k+1)<sup>3</sup>+2(k+1) is divisible by 3

(Induction)

#### Example 3

To prove  $2^n < n! \forall +ve integers n \ge 4$ .

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	n	2 <sup>n</sup>	n!	LHS < RHS?
	1	2	1	*
	2	4	2	3
	3	8	6	8
	4	16	24	
	5	32	120	
	6	64	720	<u> </u>

Prove it by induction...

(Induction)

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Induction hypothesis 2k<k!

Target: to prove  $2^{k+1} < (k+1)!$ 

> Induction step: When n=k+1,

- > LHS =  $2^{k+1}$  =  $2*2^k < 2*k! \leftarrow by hypothesis, <math>2^k < k!$
- > So, property holds for n=k+1
- > By principle of induction: property holds ∀ +ve integers n≥4

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#### Example 3

#### Why base case is n=4?

 $2^{1}=2$ When n=1, 1!=1 When n=2.  $2^2 = 4$ 2!=2  $2^3 = 8$ When n=3, 3!=6

Property does not hold for n=1, 2, 3

## Challenges ...

(Induction)

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#### Exercise

To prove  $1^2 + 2^2 + 3^2 + ... + n^2 = \frac{n(n+1)(2n+1)}{6} \forall$  +ve int  $n \ge 1$ .

n	LHS	RHS	LHS = RHS?
1	1	1*2*3/6 = 1	
2	1+4 = 5	2*3*5/6 = 5	<u> </u>
3	1+4+9 = 14	3*4*7/6 = 14	
4	1+4+9+16 = 30	4*5*9/6 =30	<u> </u>
5	1+4+9+16+25 = 55	5*6*11/6 = 55	<u> </u>

Prove it by induction...

(Induction)

Induction hypothesis: 
$$1^2 + 2^2 + 3^2 + ... + k^2 = \frac{k(k+1)(2k+1)}{6}$$

Target: to prove 
$$1^2 + 2^2 + 3^2 + ... + k^2 + (k+1)^2 = \frac{(k+1)(k+2)(2k+3)}{6}$$

Induction Step: When n = k+1

LHS = 
$$1^2 + 2^2 + 3^2 + ... + k^2 + (k+1)^2$$

← by hypothesis

#### Exercise 2

Prove that  $1+3+5+...+(2n-1) = n^2 \forall +ve integers \ge 1$ 

- > Base case: When n=1, LHS=2\*1-1=1, RHS=12=1
- > Induction hypothesis:

Assume property holds for some integer k, i.e., assume  $1+3+5+...+(2k-1)=k^2$ 

> Induction step: When n=k+1, LHS becomes 1+3+5+...+(2(k+1)-1)RHS becomes (k+1)2

Target: to prove  $1+3+5+...+(2k-1)+(2(k+1)-1))=(k+1)^2$  Exercise

To prove 
$$1^2 + 2^2 + 3^2 + ... + n^2 = \frac{n(n+1)(2n+1)}{6}$$

- > Base case: when n=1, LHS =  $1^2 = 1$ , RHS =  $\frac{1 \times 2 \times 3}{6} = 1$ =LHS
- > Induction hypothesis: Assume property holds for n=k > i.e., assume that  $1^2 + 2^2 + 3^2 + ... + k^2 = \frac{k(k+1)(2k+1)}{k}$
- > Induction step: When n=k+1, LHS becomes  $1^2 + 2^2 + 3^2 + ... + k^2 + (k+1)^2$ RHS becomes  $\frac{(k+1)(k+2)(2k+3)}{(k+1)(k+2)(2k+3)}$ Target is to prove

 $1^{2} + 2^{2} + 3^{2} + ... + k^{2} + (k+1)^{2} = \frac{(k+1)(k+2)(2k+3)}{(k+1)(k+2)(2k+3)}$ 

work on LHS & work on RHS and show that they reach the same expression

(Induction)

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Exercise 2

Prove that  $1+3+5+...+(2n-1) = n^2 \forall$  +ve integers  $\geq 1$ 

(sum of the first n odd integers equals to n2)

n	LHS	RHS	LHS = RHS?
1	1	12 = 1	
2	1+3 = 4	2 <sup>2</sup> = 4	
3	1+3+5 = 9	3 <sup>2</sup> = 9	
4	1+3+5+7 = 16	4 <sup>2</sup> =16	<u> </u>
5	1+3+5+7+9 = 25	5 <sup>2</sup> = 25	

Prove it by induction...

Induction hypothesis:  $1+3+5+...+(2k-1) = k^2$ 

Target: to prove  $1+3+5+...+(2k-1)+(2(k+1)-1))=(k+1)^2$ 

> Induction step: When n=k+1,

LHS = 1+3+5+...+(2k-1)+(2(k+1)-1)

RHS =  $(k+1)^2 = k^2 + 2k + 1 = LHS$ 

Therefore, property holds for n=k+1

By principle of induction, property holds for all +ve integers

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#### Note

The <u>induction step</u> means that if property holds for some integer k, then it also holds for k+1.

It does <u>NOT</u> mean that the property must hold for k nor for k+1.

Therefore, we  $\underline{\text{MUST}}$  prove that property holds for some starting integer  $n_0$ , which is the  $\underline{\text{base}}$   $\underline{\text{case}}$ .

Missing the base case will make the proof fail.

(Induction)

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# What's wrong with this?

More advanced ...

Claim: For all n, n=n+1

- Assume the property holds for n=k, i.e., assume k = k+1
- > Induction Step:
  - > Add 1 to both sides of the induction hypothesis
  - > We get: k+1 = (k+1)+1, i.e., k+1 = k+2
- > The property holds for n=k+1

BUT, we know this isn't true, what's wrong?

(Induction) Algorithmic Foundations COMP108 Take a student k+1 say A Induction hypothesis: same gender k В В k Swap a student, What's wrong? say B, with A Induction hypothesis: same gender

So, A, B & other (k-1) students are of the same gender

(Induction

#### What about this?

- Claim: All comp108 students are of the same gender
- > Base case: Consider any group of ONE comp108 student. Same gender, of course.
- > Induction hypothesis: Assume that any group of k comp108 students are of same gender
- Induction step: Consider any group of k+1 comp108 students...

(Induction)

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Recall: Finding minimum

Consider at the end of the statement \*\*

Base case: When i=1, M is min(a[1])

Induction hypothesis: Assume the property holds when i=k for some ke1.

Induction step: When i=k+1,

If a[k+1] < min(a[1],...,a[k]),

M is set to a[k+1], i.e., min(a[1],...,a[k+1]),

Else, a[k+1] is not min,

M is unchanged & M equals min(a[1],...,a[k+1])

Property: After each iteration of statement \*\*, the value of M is min(a[1], ..., a[i])

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