Robotics and Autonomous Systems Lecture 22: Communication in Jason

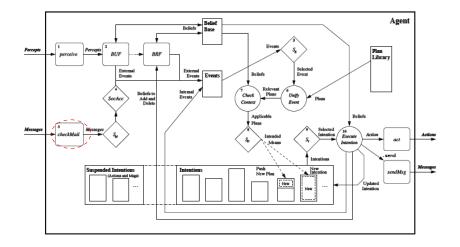
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- We will look at communication in Jason.
- This is important since you will have to write agents that communicate as part of the second assignment.
- We will look at general aspects of communication.
- We will then look at a specific example, that of the contract net.



• Each message received by the checkMail method (receiver's perspective) should be thought has having the form:

<sender, performative, content>

• Where:

- sender is the AgentSpeak term with which the agent is identified in the system
- performative this represents the goal the sender intends to achieve by sending the message
- content is an AgentSpeak formula (varying depending on the performative)

- Messages are passed through the use of internal actions that are pre-defined in Jason
- The most typical:
 - .send(receiver, performative, content)
 - .broadcast(performative, content)

where receiver, performative and content are as above

- The receiver could also be a list of agent terms
- The .broadcast action sends the message to all agents registered in the system

• The .send and .broadcast actions generate messages of the type <sender, performative, content>

which are obtained by the checkMail method of \mathbf{r} (the receiver)

• These messages (recall the previous lecture) are "filtered" during the deliberation cycle of **r** by the SocAcc function which can possibly discard them

(e.g., because of the type of sender)

- If the message goes through, Jason will interpret it according to precise semantics
 - essentially by generating new events pertaining to the goal and belief bases.
 - and **r** might then react to these events according to its plan base

tell and untell

 ${\bf s}$ intends ${\bf r}$ (not) to believe the literal in the content to be true and that ${\bf s}$ believes it

- achieve and unachieve
 s requests r (not) to try and achieve a state-of-affairs where the content of the message is true
- askOne and askAll

 ${\bf s}$ wants to know whether ${\bf r}$ knows (anybody knows) whether the content is true.

- tellHow and untellHow
 s requests r (not) to consider a plan
- askHow

 ${\bf s}$ wants to know ${\bf r}$'s applicable plan for the triggering event in the message content

• Can think of these as achieving different aims.

- tell and untell Information exchange
- achieve and unachieve Gaol Delegation
- askOne and askAll Information seeking
- tellHow and untellHow askHow
 Deliberation

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Cycle	s actions	r belief base	r events
1	<pre>.send(r, tell, open(left_door))</pre>		
2		open(left_door) [source(s)]	$\langle + open(left_door) \\ [source(s)], \top \rangle$
3	<pre>.send(r, untell, open(left_door))</pre>		
4	-		$\langle - \texttt{open(left_door)} \\ [\texttt{source(s)]}, \top \rangle$

• Information exchange

Cycle	s actions	r intentions	r events
1	<pre>.send(r, achieve, open(left_door))</pre>		
2			$\langle +! open(left_door) \\ [source(s)], \top \rangle$
3		<pre>!open(left_door) [source[s]</pre>	,
4	<pre>.send(r, unachieve, open(left_door))</pre>	<pre>!open(left_door) [source(s)]</pre>	
5	· · · ·		

Delegation

- Note that the intention is adopted after the goal is added.
- With unachieve, the internal action .drop_desire(open(left_door)) is executed.

- This semantics is operational
- Tells you how statements will be interpreted, in terms of what agents will do.
- Contrast with the mental models semantics we looked at before.

Still doesn't protect you from liars

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 .send(receiver, tellHow, "@p ... : ... <- ...") adds the plan to the plan library of r with its plan label @p • .send(receiver, untellHow, PlanLabel) removes the plan with the given plan label from the plan library of **r**

.send(receiver, askHow,

```
Goal addition event)
```

requires **r** to pass all relevant plans to the triggering event in the content (unlike for information seeking, this happens automatically)

- The CNP is a protocol for approaching distributed problem- solving
- A standardized version of the protocol has been developed by FIPA
- Agents are part of a multiagent system. They have to carry out specific tasks and they may ask other agents to perform subtasks for them
- An initiator issues a call for proposals (cfp) to all participants in the system requesting bids for performing a specific task
- After the deadline has passed, the initiator evaluates the bids it received and selects one participant to perform the task

• The contract net includes five stages:

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- Recognition;
- 2 Announcement;
- 3 Bidding;
- 4 Awarding;
- 5 Expediting.

- In this stage, an agent recognises it has a problem it wants help with.
- Agent has a goal, and either...
 - realises it cannot achieve the goal in isolation does not have capability;
 - realises it would prefer not to achieve the goal in isolation (typically because of solution quality, deadline, etc)
- As a result, it needs to involve other agents.

- In this stage, the agent with the task sends out an announcement of the task which includes a specification of the task to be achieved.
- Specification must encode:
 - description of task itself (maybe executable);
 - any constraints (e.g., deadlines, quality constraints).
 - meta-task information (e.g., "bids must be submitted by...")

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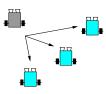
• The announcement is then broadcast.

- Agents that receive the announcement decide for themselves whether they wish to bid for the task.
- Factors:
 - agent must decide whether it is capable of expediting task;
 - agent must determine quality constraints & price information (if relevant).
- If they do choose to bid, then they submit a tender.

- Agent that sent task announcement must choose between bids & decide who to "award the contract" to.
- The result of this process is communicated to agents that submitted a bid.
- The successful contractor then expedites the task.
- May involve generating further manager-contractor relationships: sub-contracting.
 - May involve another contract net.

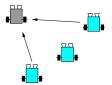
Stages

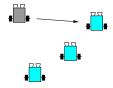




Recognition

Announcement

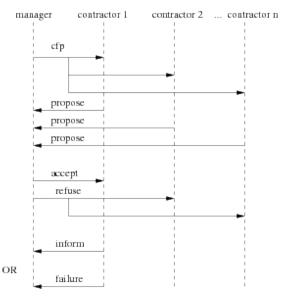




Bidding

Awarding

CNP Messages in FIPA



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MAS cnp {

```
infrastructure: Centralised
```

```
agents:
```

- c; // the CNP initiator

- pn; // a participant that does
 // not answer

}

• Here's the MAS definition

```
// Beliefs
plays(initiator,c).
```

// Plans
+plays(initiator,In)
 . .my_name(Me)
 <- .send(In,tell,introduction(participant,Me)).</pre>

// Nothing else

- Initial belief that c is the initiator.
- The belief that In is the initiator generates a message introducing itself.
- Nothing else.
- So, no response to any message

An agent that doesn't respond



```
// Beliefs
plays(initiator,c).
```

```
// Plans
+plays(initiator,In)
```

- : .my_name(Me)
- <- .send(In,tell,introduction(participant,Me)).

```
+cfp(CNPId,_Service)[source(A)] // How to respond
```

- : plays(initiator,A) // to a CfP
- <- .send(A,tell,refuse(CNPId)).

- Initial belief that c is the initiator.
- The belief that In is the initiator generates a message introducing itself.
- A CfP message from an initiator will generate a refuse message to that agent.

// Beliefs

```
plays(initiator,c).
```

price(_Service,X) :- .random(R) & X = (10*R)+100.

- Usual information about initiator
- price generates a random value for the service.

// Plans

- +plays(initiator,In)
 - : .my_name(Me)
 - <- .send(In,tell,introduction(participant,Me)).
- Usual response to finding out about the initiator.

// Plans

- Respond to CfP by making an offer.
- A proposal is added to the belief base to remember what was offered.

```
@r1 +accept_proposal(CNPId)
  : proposal(CNPId,Task,Offer)
  <- .print("My proposal '",Offer,"' won CNP ",CNPId,
            " for ",Task,"!").
@r2 +reject_proposal(CNPId)
  <- .print("I lost CNP ",CNPId, ".");
        -proposal(CNPId,_,_). // clear memory</pre>
```

- · How to handle accept and reject messages.
- Note that there is nothing here to actually do the task.
- Refusal deletes the proposal from memory.

// Beliefs

```
all_proposals_received(CNPId)
:- .count(introduction(participant,_),NP) &
    .count(propose(CNPId,_), NO) &
    .count(refuse(CNPId), NR) &
    NP = NO + NR.
```

// Goals

!startCNP(1,fix(computer)).

• all_proposals counts up the proposals received.

// Plans

Send out CfP and wait for responses

// Plans

// receive proposal @r1 +propose(CNPId,_Offer) : cnp_state(CNPId,propose) & all_proposals_received(CNPId) <- !contract(CNPId). // receive refusals @r2 +refuse(CNPId) : cnp_state(CNPId,propose) & all_proposals_received(CNPId)

- <- !contract(CNPId).
- Here we use state information.
- If every agent has responded, then go straight to awarding the contract.

```
// Needs to be atomic so as not to accept
// proposals or refusals while contracting
@lc1[atomic]
+!contract(CNPId)
   : cnp_state(CNPId,propose)
   <- -+cnp_state(CNPId,contract);
      .findall(offer(0,A),propose(CNPId,0)[source(A)],L);
      .print("Offers are ",L);
      // must make at least one offer
      L \== []:
      // sort offers, the first is the best
      .min(L,offer(WOf,WAg));
      .print("Winner is ",WAg," with ",WOf);
      !announce_result(CNPId,L,WAg);
      -+cnp state(CNPId.finished).
```

Pick an offer and announce a contract

@lc2 +!contract(_).

- We need a failure case what to do if contract is called when we aren't in the proposal state.
- Why would this happen?

// Plans

-!contract(CNPId)
 <- .print("CNP ",CNPId," has failed!").</pre>

• If the contract goal fails for some reason.

```
// Plans
```

```
+!announce_result(_,[],_).
// announce to the winner
+!announce_result(CNPId,[offer(_,WAg)|T],WAg)
    <- .send(WAg,tell,accept_proposal(CNPId));
    !announce_result(CNPId,T,WAg).
// announce to others
+!announce_result(CNPId,[offer(_,LAg)|T],WAg)
    <- .send(LAg,tell,reject_proposal(CNPId));
    !announce_result(CNPId,T,WAg).</pre>
```

- How to send out the results.
- The first clause is the base case for the recursion do nothing.



- This lecture investigated the issue of communication in Jason, highlighting the commands for the creation of messages with different performatives.
- Some of the commands were then illustrated by discussing the code of a multiagent system implementing (a stripped down version of) the contract net protocol
- As a result, the lecture also contained a brief discussion of the contract net.