

## Example: Proof by induction

For all integers  $n \geq 8$ ,  $n\text{¢}$  can be obtained using 3¢ and 5¢ coins.

**Base Case:** For  $n = 8$ ,  $8\text{¢} = 3\text{¢} + 5\text{¢}$ .

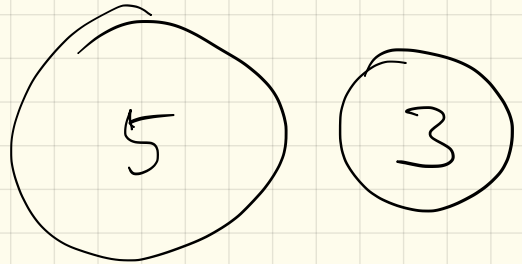
**Inductive Step:** Suppose that  $m\text{¢}$  can be obtained using 3¢ and 5¢ coins for any  $m \geq 8$ . We must show that  $(m + 1)\text{¢}$  can be obtained using 3¢ and 5¢ coins.

Consider cases

- There is a 5¢ coin among those used to make up the  $m\text{¢}$ .
  - Replace the 5¢ coin with two 3¢ coins. We obtain  $(m + 1)\text{¢}$ .
- There is no 5¢ coin among those used to make up the  $m\text{¢}$ .
  - There are three 3¢ coins ( $m \geq 8$ ).
    - Replace the three 3¢ coins with two 5¢ coins

Base case

for  $n=8$

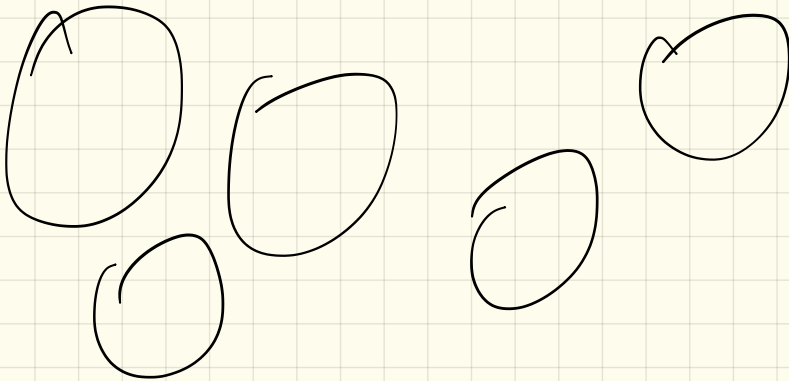


Inductive step

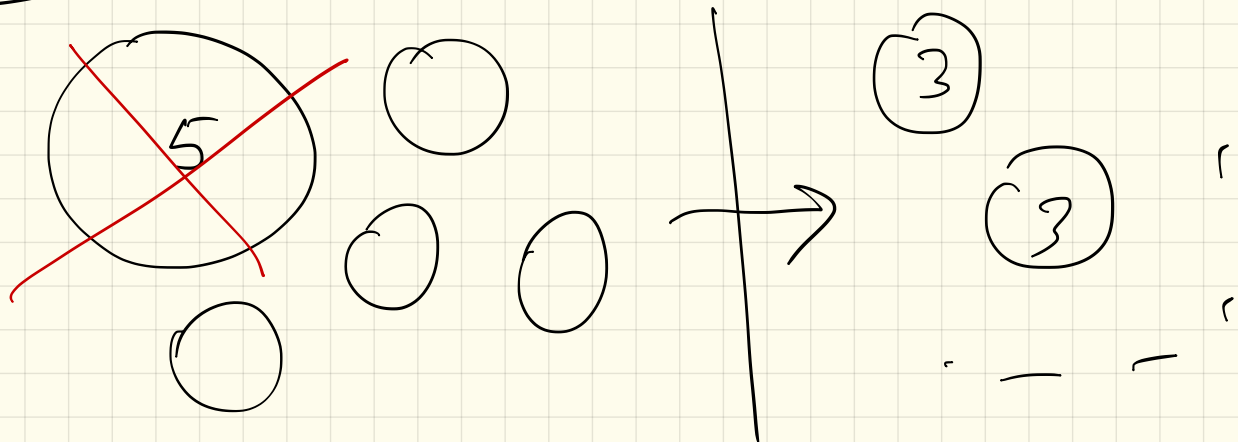
Suppose you can give  $n=m$  cents

in 3 and 5 cent coins. Then you can

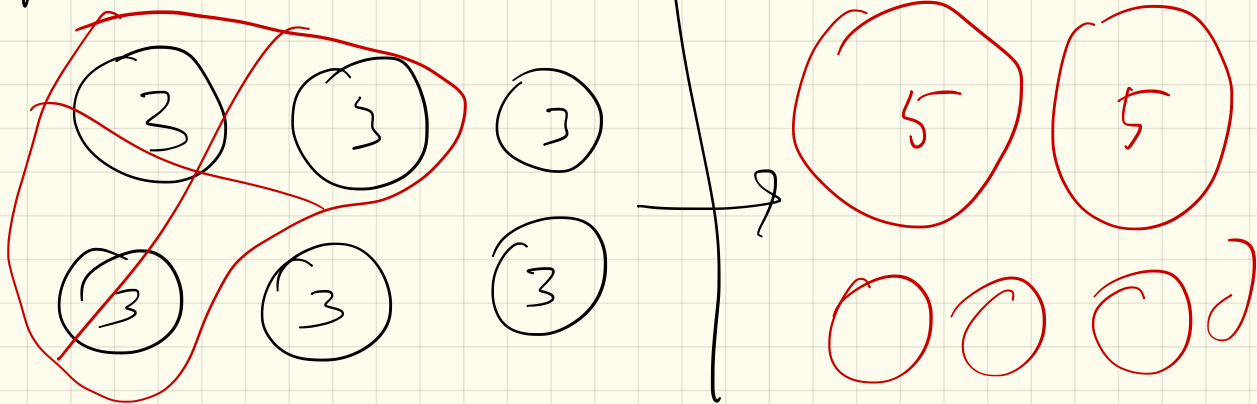
give  $n=m+1$  cents.



Case 1 There's one 5¢ coin



Case 2: No 5¢ coin on the table



## Example: Proof by induction

For every integer  $n \geq 3$ ,  $4^n > 2^{n+2}$ .

**Base Case:** Take  $n = 3$ . Then  $4^n = 4^3 = 64$ . Also,  $2^{n+2} = 2^5 = 32$ . So  $4^n > 2^{n+2}$ .

**Inductive Step:** For any  $m \geq 3$ , assume that the statement  $4^m > 2^{m+2}$  is true. (This is called the “inductive hypothesis”.) Now consider  $n = m + 1$ . We must show that  $4^{m+1} > 2^{(m+1)+2} = 2^{m+3}$ .

Here is the calculation.  $4^{m+1} = 4 \times 4^m$ . But by the inductive hypothesis,  $4 \times 4^m > 4 \times 2^{m+2}$ . Finally,

$$4 \times 2^{m+2} > 2 \times 2^{m+2} = 2^{m+3}.$$

For now define  $\forall n \geq 1$ ,  $n$  is a natural number

$$n! = 1 \times 2 \times 3 \times \dots \times n$$

Examples

$$1! = 1$$

$$2! = 1 \times 2 = 2$$

$$3! = 1 \times 2 \times 3 = 6$$

$$4! = 1 \times 2 \times 3 \times 4 = 24$$

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def fac(n):

if  $n=1$ :

return 1

else:

return  $n * \text{fac}(n-1)$

Statement:

$\forall$  natural  $n \geq 1$

fac(n) computes  $n!$

Proof We prove this statement  
by mathematical induction

Base case:  $n=1$  Then  $\text{fac}(n) = \text{fac}(1) = 1 = 1!$

Inductive step

Assume that  $\text{fac}(n) = n!$  for some  $n = m$ .

We need to show that  $\text{fac}(n) = n!$  for  $n = m+1$ .

$$\text{fac}(m) = m!$$

$\text{fac}(m+1) = (m+1) * \text{fac}(m)$ . By induction hypothesis, |

$$\begin{aligned} \text{fac}(m+1) &= (m+1) * m! = (m+1) * 1 * 2 * \dots * m = \\ & 1 * 2 * \dots * m * (m+1) = (m+1)! \end{aligned}$$

So,  $\text{fac}(n)$  computes  $n!$

# Using induction to show that a program is correct

What does the following program do?

```
    i = 0
    M = 0
    mylist = [1, 2, 6, 3, 4, 5]
    while i < len(mylist):
        M = max(M, mylist[i])
        i = i + 1
    print M
```

*my list*

→

iterations

|   | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |  |
|---|---|---|---|---|---|---|---|---|--|
| i | 0 | 1 | 2 | 3 | 4 | 5 | 6 |   |  |
| M | 0 | 1 | 2 | 6 | 6 | 6 | 6 |   |  |



## Using induction to show that a program is correct

```
i = 0
M = 0
mylist = [1, 2, 6, 3, 4, 5]
while i < len(mylist):
    M = max(M, mylist[i])
    i = i + 1
print M
```

Property: After the statement  $M = \max(M, \text{mylist}[i])$  gets executed, the value of  $M$  is  $\max(\text{mylist}[0], \dots, \text{mylist}[i])$ .

Property: After the statement  $M = \max(M, \text{mylist}[i])$  gets executed, the value of  $M$  is  $\max(\text{mylist}[0], \dots, \text{mylist}[i])$ .

**Base Case:** Take  $i=0$ . Before the statement,  $M=0$ , so the statement assigns  $M$  to be the maximum of 0 and  $\text{mylist}[0]$ , which is  $\text{mylist}[0]$ .

**Inductive Step:** Assume that the statement is true for  $i=m$  for some  $m \geq 0$ . Now consider  $i=m+1$ . The statement assigns  $M$  to be the maximum of  $\text{mylist}[m+1]$  and  $\max(\text{mylist}[0], \dots, \text{mylist}[m])$ , so after the statement,  $M$  is  $\max(\text{mylist}[0], \dots, \text{mylist}[m+1])$ .

# Strong induction

- Prove that the property holds for the natural number  $n = 0$ .
- Prove that **if** the property holds for  $n = 0, 1, 2, \dots, m$  (and not just for  $m$ !) **then** it holds for  $n = m + 1$ .

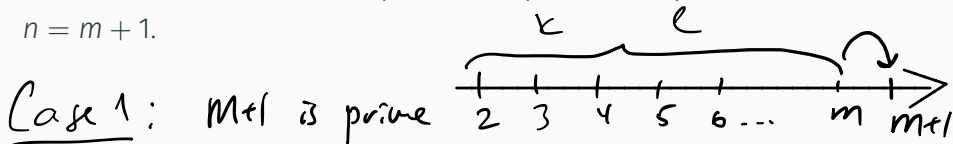
Can also be used to prove a property for all integers greater than or equal to some particular natural number  $b$

## Example: Proof by strong induction

Every natural number  $n \geq 2$ , is a prime or a product of primes.

**Base Case:** Take  $n = 2$ . Then  $n$  is a prime number.

**Inductive Step:** Assume that the property holds for  $n = m$  so every number  $i$  s.t.  $2 \leq i \leq m$  is a prime or a product of primes. Now consider  $n = m + 1$ .



There is nothing to prove

Case 2  $m+1$  is composite. Then there exist  $k, l$  natural numbers,  $k > 1, l > 1$

Then  $k$  is either prime or a product of primes

$l$  is either prime or a product of primes.

But then  $k \cdot l$  is a product of primes.

## Example: Number of multiplications

For any integer  $n \geq 1$ , if  $x_1, x_2, \dots, x_n$  are  $n$  numbers, then no matter how the parentheses are inserted into their product, the number of multiplications used to compute the product is  $n - 1$ .

$$\left( \left( \left( \left( x_1 * x_2 \right) * x_3 \right) * x_4 \right) \dots * x_n \right)$$
$$\left( \dots \right) * \left( \dots \right)$$

## Bad proofs: Arguing from example

An incorrect “proof” of the fact that the sum of any two even integers is even.

*This is true because if  $m = 14$  and  $n = 6$ , which are both even, then  $m + n = 20$ , which is also even.*

Consider the following “proof” fragment:

*Suppose  $m$  and  $n$  are any odd integers. Then by definition of odd,  $m = 2k + 1$  and  $n = 2k + 1$  for some integer  $k$ .*



## Bad proofs: Jumping to a conclusion

To jump to a conclusion means to allege the truth of something without giving an adequate reason.

*Suppose  $m$  and  $n$  are any even integers. By definition of even,  $m = 2r$  and  $n = 2s$  for some integers  $r$  and  $s$ . Then  $m + n = 2r + 2s$ . So  $m + n$  is even.*

To engage in circular reasoning means to assume what is to be proved.

*Suppose  $m$  and  $n$  are any odd integers. When any odd integers are multiplied, their product is odd. Hence  $mn$  is odd.*

## Bad proofs: Confusion between what is known and what is still to be shown

*Suppose  $m$  and  $n$  are any odd integers. We must show that  $mn$  is odd. This means that there exists an integer  $s$  such that*

$$mn = 2s + 1.$$

*Also by definition of odd, there exist integers  $a$  and  $b$  such that*

$$m = 2a + 1 \text{ and } n = 2b + 1.$$

*Then*

$$mn = (2a + 1)(2b + 1) = 2s + 1.$$

*So, since  $s$  is an integer,  $mn$  is odd by definition of odd.*

### State your game plan.

*A good proof begins by explaining the general line of reasoning, for example, “We use case analysis” or “We argue by contradiction.”*

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<sup>2</sup>*Mathematics for Computer Science* by E. Lehman, F. T. Leighton, and A. R. Meyer.

### Keep a linear flow.

*Sometimes proofs are written like mathematical mosaics, with juicy titbits of independent reasoning sprinkled throughout. This is not good. The steps of an argument should follow one another in an intelligible order.*

**A proof is an essay, not a calculation.**

*Many students initially write proofs the way they compute integrals. The result is a long sequence of expressions without explanation, making it very hard to follow. This is bad. A good proof usually looks like an essay with some equations thrown in. Use complete sentences.*

## Structure your proof

- **Theorem**—A very important true statement.
- **Proposition**—A less important but still interesting statement.
- **Lemma**—A true statement used to prove other statements.
- **Corollary**—A simple consequence of a theorem or a proposition.

## Finish

*At some point in a proof, you'll have established all the essential facts you need. Resist the temptation to quit and leave the reader to draw the “obvious” conclusion. Instead, tie everything together yourself and explain why the original claim follows.*